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Minimum number of non-zero-entries in a 7×7 stable matrix $\stackrel{\mbox{\tiny $\widehat{$}$}}{\sim}$



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ABSTRACT

We prove that if a 7×7 sign pattern matrix is potentially stable, then it has at least 11 non-zero entries. The results for $n \times n$ matrix with n up to 6 are known previously. We prove the result by making a list of possible associated digraphs with at most 10 edges, and then use algebraic conditions to show all of these digraphs or matrices cannot be potentially stable. In relation to this, we also determine the minimum number of edges in a strongly connected digraph depending on its circumference.

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1. Introduction

The concept of stability of equilibrium is central to the studies of differential equations. By using the techniques of linearization and transforming the equilibrium to zero, the

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stability problem is reduced to u' = Au, where $u \in \mathbb{R}^n$ and A is a real-valued $n \times n$ matrix. The equilibrium u = 0 is asymptotically stable if each solution u of u' = Au converges to zero as $t \to \infty$. From the theory of linear differential equation, this is equivalent to that each eigenvalue of A has negative real part. Hence it is desirable to know what kind of matrices are stable, and how to design a matrix to be stable [5].

Let M_n be the set of all $n \times n$ matrices with real-valued entries. A matrix $A \in M_n$ is said to be *stable* if, for each of its eigenvalues $\lambda_1, \lambda_2, \ldots, \lambda_n$, $\operatorname{Re}(\lambda_i) < 0$. A system which is modeled by such a matrix A has stable equilibria, and given small perturbations of its initial conditions the system will return to these equilibrium points.

We define the sign pattern of a matrix $A = [a_{ij}]$ to be an $n \times n$ matrix $S(A) = [s_{ij}]$ such that, for $i, j \in \{1, \ldots, n\}$, $s_{ij} = 0$ when $a_{ij} = 0$, $s_{ij} = -$ when $a_{ij} < 0$, and $s_{ij} = +$ when $a_{ij} > 0$. If some matrix $A \in M_n$ is found to be stable, then the sign pattern S(A)is said to be potentially stable, or PS for short. In the case where $A \in M_n$ is an upper triangular, lower triangular, or diagonal matrix, or when A is permutationally similar to such a matrix, the problem becomes trivial due to the ease of calculating the eigenvalues of these matrices. Therefore we restrict our examination to *irreducible* matrices, or the matrices $A \in M_n$ such that there does not exist a permutation matrix P such that

$$PAP^{T} = \begin{bmatrix} A_{11} & 0\\ A_{12} & A_{22} \end{bmatrix}, A_{11} \in M_{k}, A_{22} \in M_{n-k}.$$

The following result has been proved in [2].

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Theorem 1. Let the minimum number of nonzero entries required for an $n \times n$ irreducible sign pattern to be potentially stable be given by m_n . Then

$$\begin{cases} m_n = 2n - 1, & n = 2, 3, \\ m_n = 2n - 2, & n = 4, 5, \\ m_n = 2n - 3, & n = 6, \\ m_n \le 2n - (\lfloor \frac{n}{3} \rfloor + 1), & n \ge 7. \end{cases}$$

Hence the value of m_n for n = 2, 3, 4, 5, 6 was determined in Theorem 1, as well as an upper bound for m_n for any $n \ge 6$ via an explicit construction. Previously other partial results have been obtained for the cases $3 \le n \le 5$ [4,3]. In this paper we prove the following theorem:

Theorem 2. The minimum number of nonzero entries in an irreducible 7×7 real matrix is 11. Using the notation in Theorem 1, we have

$$m_7 = 2(7) - 3 = 11.$$

Note that Theorem 1 has shown that 11 is an upper bound. One can follow the proof therein to construct the following irreducible 7×7 stable matrix.

Example 3. The following 7×7 irreducible matrix

Δ	$\begin{bmatrix} -1\\ -3\\ 0\\ 0 \end{bmatrix}$	$ \begin{array}{c} 1 \\ 0 \\ 0 \\ -1 \end{array} $	1	$ \begin{array}{c} 0 \\ 0 \\ 1 \\ 0 \end{array} $	0 0 0 1	0 0 0 0	$\begin{bmatrix} 0\\0\\0\\0 \end{bmatrix}$
A –		0		0 0 0	-		$\begin{array}{c} 0\\ 0.1\\ 0 \end{array}$

has eigenvalues (approximately): $-0.3506 \pm 1.4949i$, $-0.0259 \pm 1.0025i$, $-0.0974 \pm 0.6283i$, -0.0522, and is stable.

So here in order to prove this minimum, we need only show that there cannot exist a potentially stable 7×7 sign pattern with only 10 nonzero entries. Note that if there were a potentially stable 7×7 sign pattern with fewer than 10 nonzero entries, then we would similarly be able to construct a potentially stable pattern with 10 nonzero entries by adding additional nonzero entries (with small magnitudes) to an existing potentially stable pattern. Thus, it is sufficient to prove that no potentially stable pattern with only 10 nonzero entries exists.

In order to prove that no such sign pattern exists, we first construct a list of all digraphs with 7 vertices and 10 edges which allow for correct minors (as defined by the relationship between cycles in the graph and the minors of the associated matrix in subsection 2.2). This construction is given in Section 3. Once we have constructed this list of digraphs, we will construct the associated set of nonequivalent matrix sign patterns which have correct minors. For that purpose we utilize a variant of Routh-Hurwitz stability criterion to show that none of these candidate sign patterns have a stable realization (see Section 4). From this we will conclude that the minimum number of nonzero entries must be equal to 11.

2. Preliminaries

2.1. Digraphs

We define the digraph of an $n \times n$ matrix $A = (a_{ij})$ to be a directed graph with vertex set $V_n = \{1, \ldots, n\}$, and for each $i, j \in V_n$, there exists an edge from vertex ito vertex j if and only if $a_{ij} \neq 0$. For a digraph, we define a path as an ordered set of edges such that, for some vertices $i, j, l \in \{1, \ldots, n\}$, if the m^{th} edge in the set is defined by (i, j), then the $(m + 1)^{th}$ edge is defined by (j, l). We define the *length* of a path as the number of edges in the path. If for each pair of vertices p and q, such that $p \in \{1, \ldots, n\}, q \in \{1, \ldots, n\} \setminus \{p\}$, in a given digraph there exists a path which begins at p and ends at q, we say that the digraph is *strongly connected*. It is the case that for any $A \in M_n$, A is irreducible if and only if the digraph of A is strongly connected [1]. 138

We define a *cycle* to be a path which begins and ends at the same point, and which only intersects itself at this point. We refer to a cycle of length 1 as a *loop*. Also note that a permutation similarity which swaps the i^{th} and j^{th} rows/columns of A is reflected in the digraph of A by swapping the labels of the i^{th} and j^{th} vertices of the digraph.

The *circumference* of a digraph G is defined as the length of the longest cycle present within the graph. We write this as circ(G). Note that as the circumference decreases, the minimum number of edges needed to be strongly connected increases. The following theorem gives the minimum number of edges of a strongly connected digraph G on n vertices given that circ(G) = k.

Theorem 4. Let k, n be integers such that $2 \le k \le n$ and n = a(k-1) + b for some a > 0and $0 \le b < k-1$, define $e_{n,2} = 2(n-1)$ and for k > 2, define

$$e_{n,k} = \begin{cases} ka - 1 & \text{if } b = 0, (a \ge 2) \\ ka & \text{if } b = 1 \\ ka + b & \text{if } b > 1 \end{cases}$$

If G is a strongly connected digraph with n vertices, $\operatorname{circ}(G) = k$, then $|E| \ge e_{n,k}$. Moreover, the bound is best possible, i.e., there is a graph G_0 with n vertices, $\operatorname{circ}(G_0) = k$ and $e_{n,k}$ edges.

Proof. Let G be a strongly connected digraph with vertex set $V_n = \{1, \ldots, n\}$, edge set E and $2 \leq circ(G) = k$.

Case 1: Suppose k = 2. We proceed by induction on n. When n = 2, we have $(1,2), (2,1) \in E$, and hence $|E| \ge e_{2,2} = 2$. Suppose $n \ge 3$ and any strongly connected graph $\overline{G}(V_{n-1}, \overline{E})$ with $\operatorname{circ}(G) = 2$ satisfies $|\overline{E}| \ge e_{n-1,2} = 2(n-2)$. Assume that $|E| < e_{n,2} = 2(n-1)$. Note that we can relabel the vertices so that $(n-1,n), (n,n-1) \in E$. Now,

$$E \subseteq V_n \times V_n = \underbrace{(V_{n-1} \times V_{n-1})}_{S_1} \cup \underbrace{(V_{n-2} \times \{n\})}_{S_2} \cup \underbrace{(\{n\} \times V_{n-2})}_{S_3} \cup \underbrace{\{(n-1,n), (n,n-1), (n,n)\}}_{S_4}$$

which is a disjoint union of sets. Thus,

 $|E| = |E \cap S_1| + |E \cap S_2| + |E \cap S_3| + |E \cap S_4| < 2(n-1)$ (2.1)

Since $(n-1, n), (n, n-1) \in E$, we also have $|E \cap S_4| \ge 2$ and so

$$|E \cap S_1| + |E \cap S_2| + |E \cap S_3| < 2(n-2)$$

Now, define the edge set $\overline{E} \subseteq V_{n-1} \times V_{n-1}$ as follows

$$\bar{E} = \underbrace{(E \cap S_1)}_{T_1} \cup \underbrace{\{(j, n-1) \mid (j, n) \in E \cap S_2\}}_{T_2} \cup \underbrace{\{(n-1, j) \mid (n, j) \in E \cap S_3\}}_{T_3}$$

That is, we obtain \overline{E} by removing the edges (n-1,n), (n,n-1), (n,n) from E and merging vertices n and n-1 into one vertex, labeling it as n-1. Thus,

 $|\bar{E}| \le |E \cap S_1| + |E \cap S_2| + |E \cap S_3| < 2(n-2)$

It is easy to verify that if $i, j \in V_{n-1}$ and there is a path from i to j in G, then there is a path from i to j in \overline{G} . Thus, \overline{G} is strongly connected. Also, if there is a cycle of length k in \overline{G} , then there is a cycle of length greater than or equal to k in G. Thus, circ(G) = 2. This contradicts the induction hypothesis. By mathematical induction, $|E| \ge e_{n,2}$.

Case 2: Next, assume $3 \le k \le n$ and n = a(k-1) + b. We will prove the theorem by induction on a. We start with the following base cases (i) b = 0, that is a = 2 and n = 2(k-1); (ii) b = 1, that is, a = 1 and n = k; and (iii) a = 1 and b > 1, that is n = k + b - 1.

- (i) Let n = 2(k-1). That is, a = 2 and b = 0. Then $e_{n,k} = 2k 1 = n + 1$. We are assuming G is strongly connected and $\operatorname{circ}(G) = k < n$. We can relabel the vertices so that there is a k-cycle formed by vertices $V_n V_{n-k}$, consisting of k edges. Additionally, there must an outgoing edge from each vertex $j \in V_{n-k}$. This gives us additional n-k distinct edges. Finally, there must be an outgoing edge from a vertex of $V_n V_{n-k}$ going to a vertex in V_{n-k} . Thus, $|E| \ge k + n k + 1 = n + 1 = e_{n,k}$.
- (ii) Let n = k. That is a = 1 and b = 1 and $e_{n,k} = k$. It is clear that $|E| \ge k = e_{n,k}$ since there must be an outgoing (equivalently, incoming) edge for each vertex.
- (iii) Let n = k + b 1 for some b > 1. That is a = 1 and $e_{n,k} = k + b = n + 1$. Using the same argument for n = 2(k 1), we get that $|E| \ge n + 1 = e_{n,k}$.

Assume that $a \ge 2$ when b > 0 and $a \ge 3$ when b = 0. Suppose further that any strongly connected graph $\overline{G} = (V_{n-k+1}, \overline{E})$ with $\operatorname{circ}(\overline{G}) = k$ satisfies $|\overline{E}| \ge e_{n-k+1,k}$. Suppose $|E| < e_{n,k}$. We can relabel the vertices so that $\{n-k+1,\ldots,n\}$ form a k-cycle in G, where k < n. We will define the digraph \widehat{G} with vertex set V_k and edge set $\widehat{E} = S_1 \cup S_2 \cup S_3$, where

$$S_{1} = E \cap \left(V_{n-k+1} \times V_{n-k+1} \right)$$
$$S_{2} = \left\{ (j, n-k+1) \mid (j, s) \in E \cap \left(V_{n-k} \times (V_{n} - V_{n-k+1}) \right) \right\}$$
$$S_{3} = \left\{ (n-k+1, j) \mid (j, s) \in E \cap \left((V_{n} - V_{n-k+1}) \times V_{n-k} \right) \right\}$$

That is, we remove the edges contained in the k-cycle and collapse vertices $n-k+1, \ldots, n$ into one vertex labeled by n-k+1. Then $|\hat{E}| \leq |E|-k < e_{n,k}-k = e_{n-k+1,k}$. Note that

 \hat{G} is strongly connected and $circ(\hat{G}) \leq k$. Note that from \hat{G} , we can define a strongly connected digraph \bar{G} with circ(G) = k by rearranging its edges, relocating and realigning the edges if necessary without removing or introducing a new edge. This contradicts the induction hypothesis. By mathematical induction, $|E| \geq e_{n,k}$.

For the last assertion, consider G_0 to be the digraph on V_n constructed as follows. For k = 2, let the edge set of G_0 be $E = \{(i, i + 1), (i + 1, i) \mid i = 1, ..., n - 1\}.$

For k > 2, construct a k-cycle $1 \rightarrow 2 \rightarrow \cdots \rightarrow k \rightarrow 1$; if there are at least k-1 vertices left, construct another cycle $k \rightarrow k+1 \rightarrow \cdots \rightarrow 2k-1 \rightarrow k$; if there are at least k-1vertices construct another cycle $2k-1 \rightarrow 2k \rightarrow \cdots \rightarrow 3k-2 \rightarrow 2k-1$, until we have either k-2 vertices left (when b=0) or b-1 vertices left with $1 \leq b < k-1$ vertex. For the former case or if b > 0, use a vertex in the last k-cycle and the remaining vertices to form either a k-1-cycle or a b-cycle. \Box

Suppose G is a strongly connected digraph with n vertices, edge set E, circumference k and contains m loops. Note that removing the loops does not change the strong connectivity of G. It follows from the preceding theorem that $|E| - m \ge e_{n,k}$.

2.2. Minors

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The following lemma is from elementary algebra and it is useful for better defining the properties of the characteristic polynomial of a stable matrix:

Lemma 5. Let $p(x) = x^2 + cx + d$ be a quadratic polynomial with real valued coefficients c, d. Then p has roots λ_1, λ_2 with $\operatorname{Re}(\lambda_1) < 0$ and $\operatorname{Re}(\lambda_2) < 0$ if and only if c > 0 and d > 0.

An $m \times m$ principal submatrix of A is a matrix $B = [b_{ij}] \in M_m$, $1 \le m \le n$ such that $b_{ij} = a_{v_i v_j}$ for some $v_1 < \ldots < v_m \in \{1, \ldots, n\}$. A principal minor of A is defined as the determinant of some principal submatrix $B = [b_{ij}]$ of A. We denote the $m \times m$ principal minor of A indexed by $v_1 < \ldots < v_m \in \{1, \ldots, n\}$ as $M(A)_{v_1, \ldots, v_m}$. For example,

$$A = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \implies M(A)_{23} = \det\left(\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}\right), \quad M(A)_{13} = \det\left(\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}\right).$$

There is a direct relationship between the minors of a matrix and its eigenvalues. The sum of all $k \times k$ principal minors of a matrix A is equal to the sum of all products of unique combinations of k eigenvalues of A. That is,

$$E_{k} = \sum_{1 \le v_{1} < \dots < v_{k} \le n} M(A)_{v_{1},\dots,v_{k}} = \sum_{1 \le u_{1} < \dots < u_{k} \le n} \lambda_{u_{1}} \cdots \lambda_{u_{k}}.$$
 (2.2)

Furthermore, the coefficient of t^{n-k} in the characteristic polynomial $P_A(t) = \det(tI - A)$ of the A is equal to $(-1)^k E_k$. Due to the relationship between the minors and the eigenvalues of a matrix, we have the following lemma, which is well known:

Lemma 6. If $A \in M_n(\mathbb{R})$ is stable, then the following are true:

- 1. For all k = 1, ..., n, the sign of the sum of the $k \times k$ minors of A, E_k , is $(-1)^k$.
- 2. The characteristic polynomial of A,

$$P_A(t) = \det(tI - A) = \sum_{k=0}^n (-1)^k E_k t^{n-k},$$

has all positive coefficients.

Note that the above lemma gives us a necessary condition for a given matrix A to be stable. This condition will be very important in our work. If a given sign pattern can be realized by a real valued matrix A which meets the condition that the sign of the sum of the $k \times k$ minors of A is $(-1)^k$, then we say that this sign pattern has correct minors. If for some k, the sum of $k \times k$ minors is equal to zero, then that sign pattern cannot be PS, as this would imply that either some of its eigenvalues are positive and some are negative, or that at least one of the eigenvalues is equal to zero.

The condition on the coefficients of $P_A(t)$ is necessary for the stability of A, however it is not sufficient. For example if $A = \begin{bmatrix} -0.8 & -0.81 & -1.01 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$, then

$$P_A(t) = t^3 + 0.8t^2 + 0.81t + 1.01 = (t+1)(t-0.1+i)(t-0.1-i).$$

So $P_A(t)$ has positive coefficients, but A has eigenvalues $\lambda = 0.1 \pm i$ which have strictly positive real parts, and so A is not stable.

2.3. Digraph cycles

There is a direct relationship between the minors of a matrix and the cycles present in its digraph. If two or more cycles do not share any vertices, then we say that they are *independent*. If the digraph of a sign pattern contains a cycle made up of k edges, then this implies that at least one of its $k \times k$ minors is not equal to zero. Additionally, if there exist independent cycles of length a_1, \ldots, a_m , then this implies that, if $\sum_{i=1}^m a_i \leq n$, at least one of its $(\sum_{i=1}^m a_i) \times (\sum_{i=1}^m a_i)$ minors is not equal to zero. Fig. 1 provides examples of a digraph with correct minors, and one without.

Therefore, if a given digraph has independent cycles whose lengths add up to $1, \ldots, n$, then we can assign signs to the entries of the corresponding matrix such that it has correct minors.

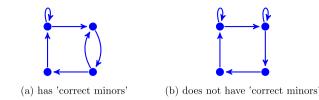


Fig. 1. A 4×4 matrix that has digraph (a) has at least one nonzero $k \times k$ minor for all k = 1, 2, 3, 4. On the other hand, a matrix with digraph (b) has all its 3×3 minors equal to zero and hence, cannot be stable.

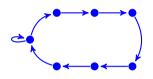


Fig. 2. Minimum configuration considered in Case 1.

3. Candidate digraphs

In this section we construct all candidate digraphs with 7 vertices and 10 edges which allow correct minors for a stable matrix. In order to better organize this list, we classify the graphs based on its circumference (the maximum length of cycle in the graph).

Case 1: circ(G) = 7.

In this case there must be at least one loop (see Fig. 2 for the minimum configuration), and either there are at least two loops or there is exactly one loop and a 2-cycle.

Case 1.1: There are at least two loops. Then 9 edges have been utilized. Suppose another edge is added to create a k-cycle where k < 7. The possible sizes of nonzero minors are 1, 2, k, k + 1, k + 2, 7 (possibly less if the k cycle intersects any of the two loops.) Thus, there is at least one 3 < r < 7 such that the $r \times r$ minor of the adjacency matrix is zero. Therefore G is not potentially stable.

Case 1.2: There is exactly one loop and a 2-cycle of two adjacent (numbering-wise) vertices. This utilizes 9 edges. Suppose the remaining edge is contained in a k-cycle, where $2 \le k < 7$. If the 2-cycle and the loop have a vertex in common, then the possible sizes of nonzero minors are 1, 2, k, k + 1, k + 2, 7, so we miss at least one minor, and therefore G is not potentially stable. Similarly, if the k-cycle has a vertex in common with either the loop or the 2-cycle, we get a non-PS adjacency matrix. Thus, the 2-cycle, loop and k cycle must be pairwise disjoint. In this case, the possible sizes of nonzero minors are 1, 2, 3, k, k + 1, k + 2, k + 3, 7. Thus, k = 4 or k = 3. In this case, we have the candidate digraphs as shown in Fig. 3.a-d.

Case 1.3: There is exactly one loop and a 2-cycle of two non-adjacent (numbering-wise) vertices. Suppose these two additional edges create a k-cycle and an r-cycle and nonzero minors of sizes 1, 2, 3, k, r, r+1, 7. Thus, (k, r) = (4, 5). In this case, we have the candidate digraph shown in Fig. 4.

Case 2: circ(G) = 6.

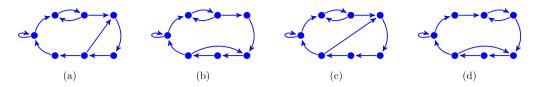


Fig. 3. Candidate digraphs under Case 1.2.

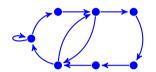


Fig. 4. Candidate digraphs under Case 1.3.

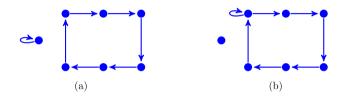


Fig. 5. Minimum configuration considered in Case 2.

In this case there must be at least one loop. Either there is a loop on the vertex that does not belong to the 6-cycle or there is none (see the two possible configurations in Fig. 5.)

Case 2.1: There is a loop in the lone vertex (say v_1) and another loop in another vertex. So far, we can guarantee nonzero minors of size 1, 2, 6, 7. For the two remaining edges, one must be outgoing from v_1 and one must be incoming from v_1 . If these two edges form a k-cycle (which intersects a loop and the 6-cycle), then we get nonzero minors of size k and k + 1 and nothing else. Thus, G will not be potentially stable.

Case 2.2: There is a loop in the lone vertex and no loop in any other vertex. Suppose the outgoing and incoming edge to the lone vertex form a k-cycle (which intersects the loop and the 6-cycle), with k < 7. Then minors of size 1, k, 6, 7 are nonzero. Suppose the remaining edge gives rise to another cycle of length 1 < r < 7 (this means it must necessarily intersect the 6-cycle. This may give rise to nonzero minors of size r, r + 1 and r + k (less if the r-cycle also intersects with the loop or the k-cycle). Thus, the r-cycle must not intersect with the k-cycle and $\{k, r, r+1, r+k\} = \{2, 3, 4, 5\}$. There is no choice but for k = 2 and r = 3. Thus, in this case, we have the candidate digraphs shown in Fig. 6.a and Fig. 6.b.

Case 2.3: There is no loop in the lone vertex. Hence, there is at least one loop intersecting the 6-cycle. Suppose the incoming and outgoing edges to the lone vertex forms a k-cycle, where 1 < k < 7. At this point, 9 edges have been accounted for and nonzero minors of sizes 1, 6, k, k + 1. Suppose the last edge gives rise to an r-cycle, where r < 7 that, by assumption, must intersect the 6-cycle. If this r-cycle intersects the loop or the

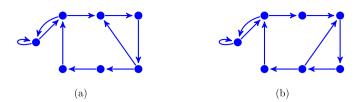


Fig. 6. Candidate digraphs under Case 2.2.

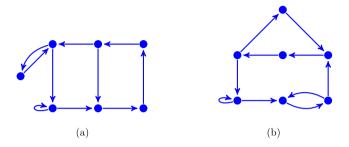


Fig. 7. Candidate digraphs under Case 2.3.

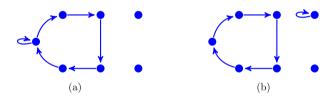


Fig. 8. Minimum configuration considered in Case 3.

k-cycle, then there will be a zero minor and thus, the adjacency matrix cannot be PS. If the *r*-cycle, loop and the *k*-cycle are pairwise disjoint, then we get additional nonzero minors of size r, r+1, r+k, r+k+1. We want $\{2, 3, 4, 5, 7\} \subseteq \{k, k+1, r, r+1, r+k, r+k+1\}$. Thus, either (r, k) = (2, 4) or (r, k) = (4, 2). Thus, in this case, we have the candidate digraphs shown in Fig. 7.a and Fig. 7.b.

Case 3: circ(G) = 5.

In this case a 5-cycle uses 5 edges, and another edge must form a loop. At least three out of the four remaining edges must be used to ensure strong connectedness of the graph. Either there is a loop intersecting the 5-cycle or there is none (see the two possible configurations in Fig. 8.)

Case 3.1: There is a loop intersecting the 5-cycle (the one we choose at the beginning, there may be more than one 5-cycle). Either there is an edge between the two remaining vertices or there is none.

Subcase 3.1.1: Suppose there is no edge connecting the two vertices (let's call them v_1 and v_2). Then, two of the four remaining edges should connect v_1 to vertices in the 5-cycle to form a k-cycle, where $k \leq 5$. Similarly, the remaining two edges must connect v_2 to vertices in the 5-cycle to form an r-cycle, where $k \leq 5$. Assuming the k-cycle, r-cycle and

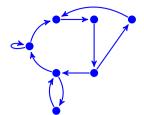


Fig. 9. Candidate digraph under Case 3.1.1.

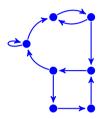


Fig. 10. Candidate digraph under Case 3.1.3.

the loop are disjoint, then we have nonzero minors of size 1, 5, k, r, k+1, r+1, k+r, k+r+1. (Note that if they are not pairwise disjoint, there will be at least minor size that will be missing.) Thus, $\{2, 3, 4, 6, 7\} \in \{k, r, k+1, r+1, k+r, k+r+1\}$. Thus, (k, r) = (2, 4) or (k, r) = (4, 2). Thus, we have the candidate digraph shown in Fig. 9.

Subcase 3.1.2: Suppose v_1 and v_2 form a 2-cycle and v_2 is not adjacent to any vertex in the 5-cycle. So far, we have accounted for 8 edges and nonzero minors of sizes 1, 2, 3, 5. The two remaining edges must be incoming and outgoing from v_1 to make a strongly connected graph. Say these two remaining edges forms a k-cycle, where $k \leq 5$. This adds nonzero minors of size k and k + 1, which is not enough to make a potentially stable adjacency matrix.

Subcase 3.1.3: Suppose v_1 and v_2 are part of a k-cycle, with $2 < k \leq 5$. So far, we have accounted for at least 9 edges and nonzero minors of sizes 1, k, 5, k+1, where ≥ 3 . To get a nonzero 2×2 minor, either there must be another loop or there is a two cycle. Adding a loop can only guarantee at least 2 more nonzero minor sizes. Thus, there must be a 2-cycle in the graph. If the two cycle is disjoint from the 5-cycle (that is, v_1 and v_2 form the 2-cycle), we only get nonzero minors of size $\{1, 2, 3, 5, k, k+1, 7\} \neq \{1, 2, 3, 4, 5, 6, 7\}$. Hence the two vertices in the 2-cycle must be part of the 5-cycle. In this case, we get nonzero minors, 1, 2, 3, k, k+1, k+2, k+3, 5. Thus, k = 4. Thus, we have the candidate digraph shown in Fig. 10.

Case 3.2: There is no loop intersecting the 5-cycle. Again, either there is an edge connecting the remaining two vertices v_1 and v_2 or there is none.

Subcase 3.2.1: There is no edge connecting the remaining two vertices v_1 and v_2 . Then, two of the four remaining edges should connect v_1 to vertices in the 5-cycle to form a k-cycle, where $k \leq 5$. Similarly, the remaining two edges must connect v_2 to vertices in the 5-cycle to form an r-cycle, where $k \leq 5$. Assuming the k-cycle, r-cycle and the loop

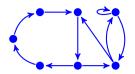


Fig. 11. Candidate digraph under Case 3.2.2.

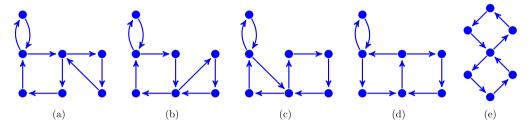


Fig. 12. Minimum configurations considered in Case 4.

are disjoint, then we have nonzero minors of size 1, 5, 6, k, r, r+1, r+k. Note that there is no choice of $2 \le k, r \le 5$ that will give a complete set of nonzero minors. Thus, this case will not give a PS adjacency matrix.

Subcase 3.2.2: Suppose v_1 and v_2 form a 2-cycle and one of v_1 or v_2 is not adjacent to any vertex in the 5-cycle. So far, we have accounted for 8 edges and nonzero minors of sizes 1, 2, 5, 6, 7. The two remaining edges must be incoming and outgoing from v_1, v_2 to make a strongly connected graph. Say these two remaining edges forms a k-cycle, where $k \leq 5$. This adds nonzero minors of size k and possibly (if the k-cycle does not contain the loop) k + 1. Thus, we have the candidate digraph shown in Fig. 11.

Subcase 3.2.3: Suppose v_1 and v_2 are part of a k-cycle, with $2 < k \leq 5$. So far, we have accounted for at least 9 edges and nonzero minors of sizes 1, 5, 6, k, where $k \geq 3$. To get a nonzero 2×2 minor, there should be another loop or a 2-cycle. Adding a loop can only guarantee at most two more sizes of nonzero minors. If there is a 2-cycle between v_1 and v_2 , we get additional nonzero minor sizes 2, 7, which is not enough for the graph to be PS. If the 2-cycle is disjoint with the k-cycle, we get additional nonzero minor sizes 2, 3, k + 2. This is still not enough to get a PS matrix.

Case 4: circ(G) = 4.

Let v_1 , v_2 , v_3 , v_4 form a 4-cycle. For each of v_5 , v_6 , v_7 , there must be incoming edges $\{5, in\}, \{6, in\}, \{7, in\}$ and outgoing edges $\{5, out\}, \{6, out\}, \{7, out\}$. Note that

$$\{\{5,in\},\{6,in\},\{7,in\}\} \cap \{\{5,out\},\{6,out\},\{7,out\}\}$$

must have at least 1 element since we still have to account for the loop. In Fig. 12, we list down all possible nonequivalent strongly connected digraphs with less than 9 edges and maximum cycle length 4.

For the digraphs in Fig. 12.a-b, adding a loop will give an adjacency matrix that has zero determinant. For the digraph in Fig. 12.c, a loop that is disjoint from the 2-cycle

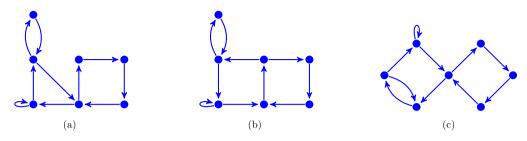


Fig. 13. Candidate digraphs under Case 4.

and 4-cycle must be added to get all nonzero minors. For the digraph in Fig. 12.d, a loop must be added so that the loop, a 4-cycle and a 2-cycle are all disjoint. Finally, for the digraph in Fig. 12.e, a loop and a 2-cycle must be added as shown in the figure below to get nonzero minors. Thus, we have the candidate digraphs shown in Fig. 13.a-c.

Case 5: circ(G) = 3.

By our formula, $e_{7,3} = 9$, so any digraph with 7 vertices and a circumference of 3 must contain at least 9 edges in order to be strongly connected. Let v_1 , v_2 and v_3 form a 3-cycle. Since the graph needs at least 9 edges in order to be strongly connected, it can have at most 1 loop, giving a total of 9 + 1 = 10 edges. Then the graph must have a 2-cycle as well.

Case 5.1: Suppose the 2-cycle shares an edge with the 3-cycle $(v_1v_2v_3)$. Then, between the 2-cycle, the 3-cycle and the loop, we have used 5 of the 10 available edges. So there are 5 edges remaining with which to connect the vertices v_4 , v_5 , v_6 and v_7 . Each of these vertices requires at least one incoming edge and one outgoing edge. Since circ(G) = 3, it would require at least 3 edges in order to connect any two of the remaining vertices to the original 3-cycle. From that point, it would require at least an additional 3 edges in order to connect the remaining two vertices. However there are only 5 edges available, and thus, the 2-cycle cannot share an edge with the 3-cycle.

Case 5.2: Suppose the 2-cycle does not share an edge with the 3-cycle $(v_1v_2v_3)$, say the 2-cycle is (v_4v_5) without loss of generality. Then, between the 2-cycle, the 3-cycle and the loop, we have used 6 of the 10 available edges. So there are 4 edges remaining with which to connect the vertices v_6 and v_7 with the cycle $(v_1v_2v_3)$ and the cycle (v_4v_5) . Since the circumference of the graph is 4, it would require at least 3 edges in order to connect v_6 and v_7 to either the 3-cycle or the 2-cycle. Then there is at most 1 edge remaining, which is insufficient to connect the remaining separated cycles. Therefore the 2-cycle cannot be separate from the 3-cycle, and so there are no digraphs with 7 vertices and a circumference of 3 which have correct minors.

Case 6: circ(G) = 2.

By our formula, $e_{7,2} = 12$, so any digraph with 7 vertices and a circumference of 2 must contain at least 12 edges in order to be strongly connected. However we are assuming only 10 edges, and therefore we cannot have a circumference of 2.

Summarizing the above discussion, we reach the main result in this section:

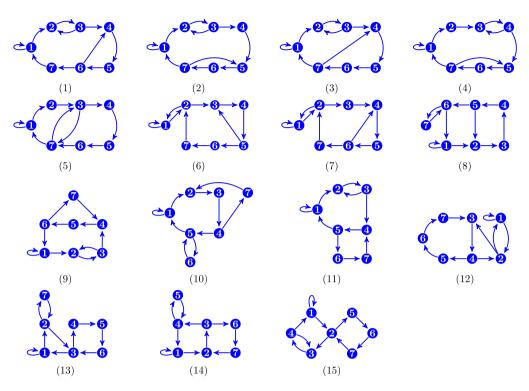


Fig. 14. List of potential digraphs with 7 vertices and 10 edges.

Proposition 7. Suppose that (V, E) is a strongly connected digraph with 7 vertices and 10 edges which has all non-zero minors. Then (V, E) is equivalent to one of digraphs in Fig. 14.

4. Calculations

We now convert the graphs from Fig. 14 into properly signed matrices, and show that none of them can be realized by a stable matrix. First however, we prove the following lemma:

Lemma 8. Let A be a 7×7 real-valued matrix with the characteristic polynomial

$$P_A(t) = t^7 + c_1 t^6 + c_2 t^5 + c_3 t^4 + c_4 t^3 + c_5 t^2 + c_6 t + c_7.$$
(4.1)

If A is stable, then all of the following inequalities must hold:

(1) $c_2c_4 - c_6 > 0;$ (2) $c_1c_2 - c_3 > 0;$ (3) $c_1c_6 - c_7 > 0;$ (4) $c_2c_5 - c_7 > 0.$

Proof. By Lemma 5, a matrix A is stable if and only if there exist $a_1, a_2, a_3, a_4, b_1, b_2, b_3 > 0$ such that

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$$P_A(t) = (t^2 + b_1 t + a_1)(t^2 + b_2 t + a_2)(t^2 + b_3 t + a_3)(t + a_4).$$
(4.2)

Comparing the coefficients of (4.1) and (4.2), we have:

$$c_{1} = a_{4} + b_{1} + b_{2} + b_{3},$$

$$c_{2} = a_{1} + a_{2} + a_{3} + a_{4}(b_{1} + b_{2} + b_{3}) + b_{1}b_{2} + b_{1}b_{3} + b_{2}b_{3},$$

$$c_{3} = (a_{1} + a_{2} + a_{3})a_{4} + (a_{2} + a_{3})b_{1} + (a_{1} + a_{3} + a_{4}b_{1})b_{2}$$

$$+ (a_{1} + a_{2} + a_{4}b_{1} + a_{4}b_{2} + b_{1}b_{2})b_{3},$$

$$c_{4} = a_{1}a_{2} + a_{1}a_{3} + a_{2}a_{3} + a_{2}a_{4}b_{1} + a_{3}a_{4}b_{1} + a_{1}a_{4}b_{2} + a_{3}a_{4}b_{2} + a_{3}b_{1}b_{2}$$

$$+ a_{1}a_{4}b_{3} + a_{2}a_{4}b_{3} + a_{2}b_{1}b_{3} + a_{1}b_{2}b_{3} + a_{4}b_{1}b_{2}b_{3},$$

$$c_{5} = (a_{1}a_{2} + a_{1}a_{3} + a_{2}a_{3})a_{4} + a_{3}(a_{2}b_{1} + a_{1}b_{2} + a_{4}b_{1}b_{2}) + a_{1}a_{2}b_{3} + a_{2}a_{4}b_{1}b_{3}$$

$$+ a_{1}a_{4}b_{2}b_{3},$$

$$c_{6} = a_{1}a_{2}a_{3} + a_{2}a_{3}a_{4}b_{1} + a_{1}a_{3}a_{4}b_{2} + a_{1}a_{2}a_{4}b_{3},$$

$$c_{7} = a_{1}a_{2}a_{3}a_{4}.$$
(4.3)

One readily sees that for each $(i, j) \in \{(1, 2), (2, 4), (1, 6), (2, 5)\}, c_i c_j$ contains all terms appearing in c_{i+j} and that $c_i c_j$ contains other terms not present in c_{i+j} .² Thus, $c_i c_j - c_{i+j}$ can be expressed as the sum of products of some of the a_k terms and b_ℓ terms. \Box

Now we use Lemma 8 to exclude all the 15 digraphs (or equivalently sign patterns) in Fig. 14 to be potentially stable. Here, we assume that the a_{ij} are positive real numbers for $1 \le i, j \le 7$.

$1. \begin{bmatrix} -a_{11} & a_{12} & 0 \\ 0 & 0 & a_{23} \\ 0 & -a_{32} & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0$	$egin{array}{ccccccccc} 0 & 0 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 \ a_{34} & 0 & 0 & 0 & 0 \ 0 & a_{45} & 0 & 0 & 0 \ 0 & 0 & a_{56} & 0 \ -a_{64} & 0 & 0 & a_{6} \ 0 & 0 & 0 & 0 \end{array}$	$c_{3} = a_{11}a_{23}a_{32} + a_{45}a_{56}a_{64}.$ So $c_{1}c_{2} - c_{3} = -a_{45}a_{56}a_{64} < 0.$ Thus this sign pattern is not potentially stable by
$2. \begin{bmatrix} -a_{11} & a_{12} & 0\\ 0 & 0 & a_{23} \\ 0 & -a_{32} & 0\\ 0 & 0 & 0\\ 0 & 0 & 0\\ 0 & 0 & 0\\ -a_{71} & 0 & 0 \end{bmatrix}$	$egin{array}{cccccccc} 0 & 0 & 0 & 0 & 0 \ 0 & 0 & 0 & 0 \ a_{34} & 0 & 0 & 0 \ 0 & a_{45} & 0 & 0 \ 0 & 0 & a_{56} & 0 \ 0 & 0 & 0 & a_6 \ 0 & -a_{75} & 0 & 0 \end{array}$	$c_{3} = a_{11}a_{23}a_{32} + a_{56}a_{67}a_{75}.$ So $c_{1}c_{2} - c_{3} = -a_{56}a_{67}a_{75} < 0.$ Thus this sign pattern is not potentially stable by

² In fact, because c_j are the *j*th elementary symmetric function of the *n* eigenvalues of *A* expressed in terms of a_k 's and b_ℓ 's in our problem, the product of the *i*th and *j*th elementary symmetric functions will contain all the terms of the (i + j)th elementary symmetric functions.

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$3. \begin{bmatrix} -a_{11} & a_{12} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & a_{23} & 0 & 0 & 0 & 0 \\ 0 & -a_{32} & 0 & a_{34} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & a_{45} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & a_{56} & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & a_{67} \\ \pm a_{71} & 0 & 0 & -a_{74} & 0 & 0 & 0 \end{bmatrix}$	Here we have: $c_2 = a_{23}a_{32}$, $c_4 = a_{45}a_{56}a_{67}a_{74}$, $c_6 = a_{23}a_{32}a_{45}a_{56}a_{67}a_{74}$. So $c_2c_4 - c_6 = 0$. Note that while both positive and negative values of a_{71} allow for correct minors, the value of a_{71} does not appear in our contradiction, and thus the contradiction holds regardless of the value of a_{71} . Thus this sign pattern is not potentially stable by Lemma 8 part 1.
$4. \begin{bmatrix} -a_{11} & a_{12} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & a_{23} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & a_{34} & 0 & 0 & 0 \\ 0 & 0 & -a_{43} & 0 & a_{45} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & a_{56} & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & a_{67} \\ -a_{71} & 0 & 0 & 0 & -a_{75} & 0 & 0 \end{bmatrix}$	Here we have: $c_1 = a_{11}$, $c_2 = a_{34}a_{43}$, $c_3 = a_{11}a_{34}a_{43} + a_{56}a_{67}a_{75}$. So $c_1c_2 - c_3 = -a_{56}a_{67}a_{75} < 0$. Thus this sign pattern is not potentially stable by Lemma 8 part 2.
$5. \begin{bmatrix} -a_{11} & a_{12} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & a_{23} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & a_{34} & 0 & 0 & a_{37} \\ 0 & 0 & 0 & 0 & a_{45} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & a_{56} & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & a_{67} \\ -a_{71} & 0 & -a_{73} & 0 & 0 & 0 & 0 \end{bmatrix}$	Here we have: $c_1 = a_{11}, c_2 = a_{37}a_{73},$ $c_3 = a_{11}a_{37}a_{73}.$ So $c_1c_2 - c_3 = 0.$ Thus this sign pattern is not potentially stable by Lemma 8 part 2.
$6. \begin{bmatrix} -a_{11} & a_{12} & 0 & 0 & 0 & 0 & 0 \\ -a_{21} & 0 & a_{23} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & a_{34} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & a_{45} & 0 & 0 \\ 0 & 0 & -a_{53} & 0 & 0 & a_{56} & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & a_{67} \\ 0 & -a_{72} & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$	Here we have: $c_1 = a_{11}, c_6 = a_{23}a_{43}a_{45}a_{56}a_{67}a_{72},$ $c_7 = a_{11}a_{23}a_{43}a_{45}a_{56}a_{67}a_{72}.$ So $c_1c_6 - c_7 = 0.$ Thus this sign pattern is not potentially stable by Lemma 8 part 3.
$7. \begin{bmatrix} -a_{11} & a_{12} & 0 & 0 & 0 & 0 & 0 \\ -a_{21} & 0 & a_{23} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & a_{34} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & a_{45} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & a_{56} & 0 \\ 0 & 0 & 0 & -a_{64} & 0 & 0 & a_{67} \\ 0 & -a_{72} & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$	Here we have: $c_1 = a_{11}, c_6 = a_{23}a_{43}a_{45}a_{56}a_{67}a_{72},$ $c_7 = a_{11}a_{23}a_{43}a_{45}a_{56}a_{67}a_{72}.$ So $c_1c_6 - c_7 = 0.$ Thus this sign pattern is not potentially stable by Lemma 8 part 3.
$8. \begin{bmatrix} -a_{11} & a_{12} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & a_{23} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & a_{34} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & a_{45} & 0 & 0 \\ 0 & -a_{52} & 0 & 0 & 0 & a_{56} & 0 \\ -a_{61} & 0 & 0 & 0 & 0 & 0 & a_{67} \\ 0 & 0 & 0 & 0 & 0 & -a_{76} & 0 \end{bmatrix}$	Here we have: $c_2 = a_{67}a_{76}, c_4 = a_{23}a_{34}a_{45}a_{52},$ $c_6 = a_{67}a_{76}a_{23}a_{34}a_{45}a_{52}.$ So $c_2c_4 - c_6 = 0.$ Thus this sign pattern is not potentially stable by Lemma 8 part 2.

$9. \begin{bmatrix} -a_{11} & a_{12} \\ 0 & 0 \\ 0 & -a_3 \\ 0 & 0 \\ 0 & 0 \\ \pm a_{61} & 0 \\ 0 & 0 \end{bmatrix}$	$a_{23} 0 0$	$ \begin{bmatrix} 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ a_{56} & 0 \\ 0 & a_{67} \\ 0 & 0 \end{bmatrix} $	Here we have: $c_2 = a_{23}a_{32}, c_5 = a_{11}a_{45}a_{56}a_{67}a_{74},$ $c_7 = a_{23}a_{32}a_{11}a_{45}a_{56}a_{67}a_{74}.$ So $c_2c_5 - c_7 = 0.$ Thus this sign pattern is not potentially stable by Lemma 8 part 4 (regardless of the sign of the (6, 1) entry).
$10. \begin{bmatrix} -a_{11} & a_{12} \\ 0 & 0 \\ 0 & 0 \\ \pm a_{51} & 0 \\ 0 & 0 \\ 0 & -a_7 \end{bmatrix}$	$egin{array}{ccccccccc} a_{23} & 0 & 0 \ 0 & a_{34} & 0 \ 0 & 0 & a_{45} \ 0 & 0 & 0 \ 0 & 0 & -a_{65} \end{array}$	$\begin{bmatrix} 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & a_{47} \\ a_{56} & 0 \\ 0 & 0 \\ 0 & 0 \\ \end{bmatrix}$	Here we have: $c_2 = a_{56}a_{65}$, $c_4 = a_{23}a_{34}a_{47}a_{72}$, $c_6 = a_{56}a_{65}a_{23}a_{34}a_{47}a_{72}$. So $c_2c_4 - c_6 = 0$. Thus this sign pattern is not potentially stable by Lemma 8 part 2 (regardless of the sign of the (5, 1) entry).
111. $\begin{bmatrix} -a_{11} & a_{12} \\ 0 & 0 \\ 0 & -a_3 \\ 0 & 0 \\ \pm a_{51} & 0 \\ 0 & 0 \\ 0 & 0 \end{bmatrix}$	$a_{23} 0 0$	$\begin{bmatrix} 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ a_{56} & 0 \\ 0 & a_{67} \\ 0 & 0 \end{bmatrix}$	Here we have: $c_2 = a_{23}a_{32}$, $c_4 = a_{45}a_{56}a_{67}a_{74}$, $c_6 = a_{23}a_{32}a_{45}a_{56}a_{67}a_{74}$. So $c_2c_4 - c_6 = 0$. Thus this sign pattern is not potentially stable by Lemma 8 part 2 (regardless of the sign of the (5, 1) entry.
$12. \begin{bmatrix} -a_{11} & a_{12} \\ -a_{21} & 0 \\ 0 & 0 \\ 0 & -a_4 \\ 0 & 0 \\ 0 & 0 \\ 0 & 0 \end{bmatrix}$	$\begin{array}{cccc} a_{23} & 0 & 0 \\ 0 & a_{24} & 0 \end{array}$	$ \begin{bmatrix} 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ a_{56} & 0 \\ 0 & a_{67} \\ 0 & 0 \end{bmatrix} $	Here we have: $c_2 = a_{12}a_{21}, c_5 = a_{34}a_{45}a_{56}a_{67}a_{73},$ $c_7 = a_{12}a_{21}a_{34}a_{45}a_{56}a_{67}a_{73}.$ So $c_2c_5 - c_7 = 0.$ Thus this sign pattern is not potentially stable by Lemma 8 part 4.
$13. \begin{bmatrix} -a_{11} & a_{12} \\ 0 & 0 \\ \pm a_{31} & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & -a_7 \end{bmatrix}$	$egin{array}{cccccccc} 0 & 0 & 0 \ a_{23} & 0 & 0 \ 0 & a_{34} & 0 \ 0 & 0 & a_{45} \ 0 & 0 & 0 \ -a_{63} & 0 & 0 \ _2 & 0 & 0 & 0 \end{array}$	$\begin{bmatrix} 0 & 0 \\ 0 & -a_{27} \\ 0 & 0 \\ 0 & 0 \\ a_{56} & 0 \\ 0 & 0 \\ 0 & 0 \end{bmatrix}$	Here we have: $c_2 = a_{27}a_{72}$, $c_4 = a_{34}a_{45}a_{56}a_{63}$, $c_6 = a_{27}a_{72}a_{34}a_{45}a_{56}a_{63}$. So $c_2c_4 - c_6 = 0$. Thus this sign pattern is not potentially stable by Lemma 8 part 2 (regardless of the sign of the (3, 1) entry).
$14. \begin{bmatrix} -a_{11} & a_{12} \\ 0 & 0 \\ 0 & 0 \\ \pm a_{41} & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & -a_{7} \end{bmatrix}$	$egin{array}{cccccccc} 0 & 0 & 0 & \ a_{23} & 0 & 0 & \ 0 & a_{34} & 0 & \ 0 & 0 & a_{45} & \ 0 & -a_{54} & 0 & \ 0 & 0 & 0 & \ 2 & 0 & 0 & 0 & \ \end{array}$	$ \begin{bmatrix} 0 & 0 \\ 0 & 0 \\ a_{36} & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & a_{67} \\ 0 & 0 \end{bmatrix} $	Here we have: $c_2 = a_{45}a_{54}, c_5 = a_{11}a_{23}a_{36}a_{67}a_{72},$ $c_7 = a_{45}a_{54}a_{11}a_{23}a_{36}a_{67}a_{72}.$ So $c_2c_5 - c_7 = 0.$ Thus this sign pattern is not potentially stable by Lemma 8 part 4 (regardless of the sign of the (4, 1) entry).

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